4SECURail

The 4SECURail Formal Methods Demonstrator

RSSRAIL 2022, June 2, Paris

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The railway infrastructure is a complex System of Systems

Spreading across many national borders

Managed by many administrative bodies

Developed by many producers

Expensive to develop, maintain and exercise safely





The solution: *High Quality Standard Interfaces between components*

- + to reduce costs and vendors lock-in
- + to increase competitivity, dependability and efficiency (safety is already guaranteed)

Several initiatives try to advance the state of art (e.g. EULYNX / ERTMS / SHIFT2RAIL / Europe's Rail)

recognizing the importance of formal analysis

(during development and during standardization)









4SECURail (November 2019 - Novemer 2021) is a (small) project of the <u>Shitf2Rail</u> initiative

One of its golas is a **controlled experiment** (*demonstrator*) in exploiting formal methods *in the requirements definition phase* of a railway signalling system.

- Can formal methods help improving the quality of requirement specifications (standards)? **How**?

- Can their adoption be cost effective for IM? **How much**?

I.e. The project takes the point of view of the Infrastructure Manager (standardization bodies), with focus not just in safety but also interoperability





The 4SECURail approach (incremental/iterative)





The 4SECURail case study

(RBC-RBC(Radio Block Centre) communications)







4SECURail: The Artifacts of the Demonstrator





Why an Executable UML/SysML model?

- Removing ambiguity in the initial NL documentation by adopting a standard, widely known, precise notation.
- Allowing, not formal methods experts, to understand and confirm the underlying design being modelled.
- Remaining at this level independent from the specific formal verification framework(s) adopted (preferrable in the case of international standards)





- FIFO events queues
- No priority conflicts
- No parallel or composite states
- No deferred events
- No history/deep-history states
- Basic data types (enum, int, bool, vectors)
- Basic statements (assignments, conditionals)
- No entry/exit/do activities





4SECURail: Executable UML Modelling (example)





4SECURail: from Executable to Formal







Formal Modelling and Analysis (1)







4SEURail: Why three formal models (UMC, ProB, LNT)?

- The three formal models can be compared for equivalence, detecting possible errors made in the formal encoding.
- The three different verification frameworks provide different verification functionalities. (e.g. linear vs branching time, compositional vs explicit)
- When the same fuctionality is supported (e.g. animation, analisys of counter examples), the most user-friendly framework can be used.
- It is however more expensive and difficult to become expert users of several verification frameworks.





4SECURail: Formal Modelling and Analysis

ProB

- Static Analysis
- Reachability Properties
- Statespace Stats
- State Invariants
- Deadlocks
- LTLe Model Checking
- CTLe Model Checking

UMC

- Static Analysis
- Reachability Properties
- System Traces Minimization
- Statespace Stats
- Deadlocks
- Runtime Errors
- UCTL Model Checking (state/event based)
- Custom system observations
- Explanations as Message
 - Sequence Diagrams

LNT

- Static Analysis
- Reachability Properties
- Statespace Stats
- Deadlocks

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- MCL Model Checking (event based)
- Compositional Verification
- Strong/ Divbranching/ Sharp Minimizations
- Powerful scripting language



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4SECURail: different levels of complexity of analysis

 Simple «push-button» like formal analysis (static analysis, reachability analysis, deadlock checking)

- More advanced verifications (model checking temporal logic formalas, compositional analysis, bisimulations and equivalences)





4SECURail: back from Formal Models to Natural Laguage







4SECURail: hiding non essential implementation details



4SECURail: The Approach of the Demonstrator





4SECURAIL: from Abstract Modelling to NL Requirements

- UML transitions directly mapped to NL requirements on control flow.
- Explicit definition, for ech component, of the assumptions it makes on the rest of the system, and the guarantees of which it is responsible.
- Rigorous specification of the syntactic interface between component.





Conclusions and Observations:

- The construction of the **executable model** already reveals all the NL **ambiguities**, *part* of the **inconsistences**, and **missing points**.
- Formal methods diversity allows to detect errors in the formal models encoding, as well as in the translation and verification tools.
- Formal analisys of the executable model allows to detect errors in the implementation, to identify hidden assumptions, and to assess the expected guarantees of the various components.
- In the really "early" stages of requirements definition, makes sense to investigate the "reverse" flow: from Formal Models to Natural Language





4SECURail website: <u>https://4securail.eu</u>

4SECURail Deliverables doi: <u>10.5281/zenodo.5807738</u>

- D2.1 Initial rationale for demonstrator structure
- D2.3 Initial case study requirements definition
- D2.5 The formal methods demonstrator experiment

Revised case study requirementsdoi: 10.5281/zenodo.5541217Formal models and scenariosdoi: 10.5281/zenodo.5541307Model transformation toolsdoi: 10.5281/zenodo.5541350





4SECURail: Structured Natural Language Requirements

Configuration Parameters .. External Interactions ... External Guarantees ... External Assumptions ...

Behavioral Requirements ...

- **R2**: When in <u>Disconnected</u> state, the CSL immediately sends a SAI_CONNECT.request to the SAI component, starts a connTimer, and moves to the Connecting state.
- **R3**: When in <u>Connecting</u> state the connTimer expires, the CSL moves to Disconnected state.





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Thanks!

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